全息纠缠熵的比特流描述及 其可能应用



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Outline

Holographic entanglement entropy and tensor network: A review

≻Quantum bit threads of MERA in large c

- ► Applications to HEoP
- ≻ Summary

Black hole entropy

From the first law we see BH has entropy, the Bekenstein-Hawking

entropy(Bekenstein 1973)

$$S = \frac{Ac^3}{4G\hbar}$$



Usually entropy is the logarithm of phase space volume, but BH entropy scales like area





Entanglement?

The black hole interior is entangled with its exterior radiations. This naturally give the area scaling of the entropy (Sorkin; Bombelli,Koul,Lee&Sorki n)







(Ryu & Takayanagi '06)



Holography & tensor network

Recently the relation between **Tensor Network of wave function in quantum critical phase** and discrete **Anti de Sitter (AdS) space** has been suggested. (B.Swingle '2009)



Why tensor network?



"The fundamental laws necessary for the mathematical treatment of a large part of physics and the whole of chemistry are thus completely known, and the difficulty lies only in the fact that application of these laws leads to equations that are too **complex** to be solved."

—Paul Dirac

$$i\hbar \frac{\partial}{\partial t} \Psi = \widehat{H} \Psi$$

Schrodinger equation





For many body system, only a tiny part of the Hilbert space is relavant. Many d.o.f. are irrelavant to the ground state in question.







For a ground state



MERA

G.Vidal PRL, 2007, 99: 220405, PRL, 2008, 101: 110501.

Multi-scale entanglement renormalization ansatz

(MERA)



MERA as AdS/CFT



 $S(A)\approx \log L$

B. Swingle, PRD, 2012, 86(6): 065007.B. Swingle, arXiv:1209.3304



MERA as AdS/CFT



B. Swingle, PRD, 2012, 86(6): 065007.B. Swingle, arXiv:1209.3304



 $S_A = \frac{Area(\gamma_A)}{AG^{d+2}} \propto \log L$

Nozaki *et al* JHEP10(2012)193

W.C. Gan(甘文聪),FWS, M.H.Wu(吴孟和), PLB 760 (2016) 796, 772 (2017) 464-470

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Bit threads

M. Freedman and M. Headrick, Commun.Math. Phys. 352, 407 (2017)

Subtlety: A strangely discontinuous transition of the bulk minimal surface under continuous deformations of *A*.



The bits encoding the microstate of A somehow "live on" the minimal surface and strangely jump from one place to another!

Max-Flow/Min-Cut





The area of minimal surface can be replaced by a maximal flux. Information is encoded in the flow, not the surface.

Bits information is quantum, but network is classical

Flow picture implies threads live in a Lorentzian manifold



C-B Chen(陈崇斌), FWS, M-H Wu(吴孟和), arXiv:1804.00441

• A classical network is incomplete for quantum bits!



• AdS time slice is a spacelike hypersurface and don't have causal structure



Quantum Max-flow/Min-cut



For fixing basic of inputs S and outputs T, each tensor assignments determine a state of this tensor network:

$$\begin{aligned} |\alpha(G,a;\mathcal{T})\rangle &:= \sum_{\substack{i_1,\cdots,i_{|S|}\\j_1,\cdots,j_{|T|}}} C_{i_1,\cdots,i_{|S|},j_1,\cdots,j_{|T|}} \\ &\times |i_1,\cdots,i_{|S|}\rangle_S |j_1,\cdots,j_{|T|}\rangle_T, \end{aligned}$$

S. X. Cui, M. H. Freedman et.al, J.Math. Phys. 57 062206 (2016)

Definition 1 (Quantum Min-cut). The quantum min-cut QMC(G, a) is the minimum value of product of capacities over all edge cut sets, i.e.

$$QMC(G, a) := \min_{A} \prod_{e \in C} a_e.$$

There is a linear map $\beta(G, a; \mathcal{T}) \in V_S^* \bigotimes V_T = Hom(V_S, V_T)$ from inputs to outputs: $V_S \mapsto V_T$ acting on the inputs state:

$$\beta(G,a;\mathcal{T})|i_1,\cdots,i_{|S|}\rangle_S:=\sum_{j_1,\cdots,j_{|T|}}C_{i_1,\cdots,i_{|S|},j_1,\cdots,j_{|T|}}\times|j_1,\cdots,j_{|T|}\rangle_T.$$

Briefly $\beta(G, a; \mathcal{T})$ is the map of big tensor consist of tensors from S to T.

Definition 2 (Quantum Max-flow). For over all tensor assignments, there exists a maximal value of the rank of map $\beta(G, a; \mathcal{T})$ and we define this maximal value as the quantum max-flow:

$$QMF(G, a) := \max_{\mathcal{T}} rank(\beta(G, a; \mathcal{T})).$$



Kinematic space

B. Czech et al JHEP 1607, 100(2016)

Kinematic space is defined as a collection of geodesics in AdS time slice



 $\omega(\theta, \alpha)$; the measure of the Kinematic space, determining the structure of the space





Kinematic space is a de Sitter space! A Lorentzian manifold with causal structure.

Kinematic space VS MERA

For causal structure consideration, kinematic space should be viewed as the corresponding geometry of the MERA

The conditional mutual information can be obtained by counting the number of edges which is the net reduction of edges through a causal diamond from bottom up

$$I(A, C|B) = S(AB) + S(BC) - S(B) - S(ABC)$$
$$\approx \frac{\partial^2 S(u, v)}{\partial u \partial v} du dv$$

The volume of a Kinematic space is determined by # of isometries

Metric of MERA is given by the one of Kinematic:

B. Czech et al JHEP 1607, 100(2016)



 $\mathcal{D}(\text{isometry}) = I(A, C|B)$

$$ds_{MERA}^2 = I(\Delta u, \Delta v | B) \xrightarrow{MERA} (\# \text{ of isometry}) \Delta u \Delta v$$



C-B Chen(陈崇斌), FWS, M-H Wu(吴孟和), arXiv:1804.00441

Now question is if QMF=QMC?

In generaly, QMF/QMC theorem cannot be satisfied, instead

 $QMF(G, a) \le QMC(G, a)$

which is different from the classical case. This implies some quantum effects.

When QMF=QMC?

(i) If the type of the tensors is not constraint, then we have For a given graph G(V,E), if the capacity a of each edge is a power of d, where d is an positive integer, then QMF=QMC.

(ii) If the type of the tensors is fixed, then: QMF=QMC is "asymptotically" true in the large $c(\log \chi \simeq \frac{c}{6})$ limit $QMF(G, \chi, 0) = QMC(G, \chi) \cdot (1 - 0(1))$ M B Hastings arXiv:

M. B. Hastings,arXiv:1603.03717

Entanglement entropy flow in MERA tensor network



C-B Chen(陈崇斌), FWS, M-H Wu(吴孟和), arXiv:1804.00441

There are $\log \chi$ (bits) entanglement entropy flow in each edge of tensor network.

From bottom up, $2 \log \chi$ (bits) entropy flow become $\log \chi$ (bits) entropy flow

By introducing a density of the isometry for flow (bits) conservation:

$$|
ho| \leq
ho_M,$$
 (Finite)
 $abla_\mu f^\mu = -
ho$ (conservation)

 $\boldsymbol{\rho} {:} \mbox{ density of isometries, the sources of sinks of the flow }$

HEE from quantum bit threads

- > Subadditivity $S(A) + S(B) \ge \int_{C_A} f(AB) + \int_{C_B} f(AB)$ $\ge \int_{C_{AB}} f(AB) = S(AB).$
- Strong subadditivity

$$\begin{split} I(A:C|B) &= \int_{C_{AB}} f(A,B,C) + \int_{C_{BC}} f(A,B,C) \\ &- \int_{C_B} f(A,B,C) - \int_{C_{ABC}} f(A,B,C) \\ &= S(AB) + S(BC) - S(B) - S(ABC) \\ &= - \int_D \rho \geq 0. \end{split}$$

> Araki-Lieb inequality

$$\begin{split} S(AB) + S(A) &\geq \int_{AB} f(B, A', A) + \int_{D_{AB}} \rho \\ &+ \int_{A'} f(B, A', A) + \int_{D_{A'}} \rho \\ &\geq \int_{A'AB} f(B, A', A) + \int_{D_{A'AB}} \rho \\ &\geq \int_{D_B} \rho + \int_B f(B, A', A) = S(B) \end{split}$$

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Entanglement of purification

- S_A does not measure entanglement when ρ_{AB} is a mixed state
- Long distance quantum teleportation or global quantum key distribution need to distribute a certain supply of pairs of particles in a maximally entangled state to two distant users

Entanglement of purification

B. M. Terhal et al, J. Math. Phys. 43 (2002) 4286

н.

Let
$$|\psi\rangle \in H_{AA'} \otimes H_{BB'}$$
 be a purification of ρ_{AB} , so that
 $\operatorname{Tr}_{A'B'} |\psi\rangle \langle \psi| = \rho_{AB}$, the entanglement of purification (EoP)
is
 $\operatorname{E}_{P}(A:B) = \min_{\psi,A'B'} S_{AA'}$
where $S_{AA'}$ is Von Neumann entropy.
 $|\psi\rangle_{AA'BB'}$
 $|\psi\rangle_{AA'BB'}$
 A
 B

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Holographic entanglement of purification

◆ Surface/state correspondence





Trivial State $|\Omega\rangle$

M. Miyaji *et al*, PTEP 2015(2015)7, 073B03, PRL 115, no.17, 171602 (2015)





Mixed State $\rho(\Sigma)$

Mixed State $\rho(\Sigma)$

◆ In AdS/CFT, the HEoP is conjectured to be given by

$$E_P(A:B) = \frac{\operatorname{area}(\sigma_{AB}^{\min})}{4G_N} \equiv E_W(A:B)$$

where σ_{AB}^{min} is the minimal cross section on the entanglement wedge r_{AB} .

K. Umemoto and T. Takayanagi, Nature Phys. 14, no. 6, 573 (2018)



Bit thread formulation of HEoP

• Define a flow from A to its complement, a vector field on the entanglement wedge r_{AB} , satisifying

$$\nabla \cdot \upsilon_{\mathrm{AB}} = 0, \left| \upsilon_{\mathrm{AB}} \right| \le \frac{1}{4G_N},$$

$$v_{AB} = -v_{AB}, \hat{\mathbf{n}} \cdot v_{AB} = 0 \text{ on } m_{AB}$$

• We asume that the EoP can be written as the flux of a max flow from A to B

$$E_{P}(A:B) = \max_{\substack{\upsilon_{AB}:\\\hat{n}\cdot\upsilon_{AB}\mid_{m_{AB}}=0}} \int_{A} \upsilon_{AB}$$

D-H Du(杜东辉), C-B Chen (陈崇斌), FWS, 1904.06871, JHEP 08(2019)140



M. Ghodrati, X.-M. Kuang, B. Wang, C.-Y. Chen, Y.-T. Zhou, arXiv:1902.02475

◆ Using the generalized Riemannian MFMC theorem

$$\max_{\substack{v_{AB}:\\\hat{n}\cdot v_{AB}|m_{AB}=0}} \int_{A} v_{AB} = \min_{\substack{\sigma_{AB}\sim A\\ rel \ m_{AB} \ on \ \partial r_{AB}}} \frac{\operatorname{area}(\sigma_{AB})}{4G_N} \equiv E_W(A:B)$$

M. Headrick and V. E. Hubeny, Class. Quant. Grav. 35, no. 10, 10 (2018)

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Bit thread formulation of the quantum advantage of dense code (QAoDC)



Choose a flow on r_{AB} that simultaneously maximizes the flux from A to (BC') and the flux from A to B, satisfying

$$\nabla \cdot \widetilde{\mathcal{O}}_{A(B,C')} = 0, \left| \widetilde{\mathcal{O}}_{A(B,C')} \right| \le \frac{1}{4G_N}, \widetilde{\mathcal{O}}_{A(B,C')} = -\widetilde{\mathcal{O}}_{(B,C')A}$$

For this flow configuration, we have

 $\tilde{v}_{A(B,C')}$

$$\mathbf{S}(\mathbf{A}) = \int_{\mathbf{A}} \widetilde{\mathcal{O}}_{A(B,C')} = \int_{B} \widetilde{\mathcal{O}}_{(B,C')A} + \int_{C'} \widetilde{\mathcal{O}}_{(B,C')A} = E_{P}(A:B) + \int_{C'} \widetilde{\mathcal{O}}_{(B,C')A}$$

• Recalling that for pure tripartite state, we have $S(A) = E_P(A:B) + \Delta(C' > A)^1$

The QAoDC can be written as $\Delta(B > A) \equiv S(A) - inf_{\Lambda_B}S[(I_A \otimes \Lambda_B)\rho_{AB}] = sup_{\Lambda_B}I'(B \land A)$

$$\Delta(C' > A) = \int_{C'} \widetilde{\nu}_{(B,C')A}$$

1. M. Horodecki and M. Piani, J. Phys. A: Math. Theor. 45 (2012) 105306

Flow-based proofs of the properties of EoP

(i) The E_P is bounded above by the entanglement entropy

 $\mathbf{E}_{\mathbf{P}}(A:B) \le \min(S(A), S(B))$

(ii) The E_P is monotonic

 $\mathbf{E}_{\mathbf{P}}(A:B\mathbf{C}) \ge \mathbf{E}_{\mathbf{P}}(A:B)$

(iii) The E_P is bounded below by half the mutual information

$$\mathrm{E}_{\mathrm{P}}(A:B) \ge \frac{\mathrm{I}(A:B)}{2}$$

(iv) The E_P is polygamous for a tripartite pure state

$$\mathbf{E}_{\mathbf{P}}(A:B) + \mathbf{E}_{\mathbf{P}}(A:C) \ge \mathbf{E}_{\mathbf{P}}(A:BC)$$

(v) For a tripartite system,

$$E_{p}(A:BC) \ge \frac{I(A:B)}{2} + \frac{I(A:C)}{2}$$

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And more

(i) Flow based proof of the monogamy relation of QAoDC with the EoP for the tripartite state:

 $\mathbf{S}(\mathbf{A}) \ge E_P(A:B) + \Delta(C > A)$

(ii) A new lower bound for S(AB) in terms of QAoDC, which is tighter than the one given by the Araki-Lieb inequality: $|S(A) - S(B)| \le S(AB)$

 $\Delta(C > A) + \Delta(C > B) \le S(AB)$

(iii) A new inequality for EoP: $E_P(A:BC) \le E_P(B:AC) + E_P(C:AB)$

Summary

- Tensor network provide a possible perspective on the relation between holographic space and entanglement
- HEE admits a bit-thread interpretation
- Classical bit-thread picture is incomplete, quantum version should be introduced.
- Bit-thread interpretation can be applied to some fields, such as holographic entanglement of purification.

Thank you

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